

## PS1 - Q1

For this recurrence, we will first multiply both sides by  $N$ :

$$NA_N = N + 2 \sum_{1 \leq j \leq N} A_{j-1} \quad \text{for } N > 0 \quad (1)$$

Next subtract the same equation for  $N - 1$ :

$$NA_N - (N - 1)A_{N-1} = 1 + 2A_{N-1} \quad \text{for } N > 1 \quad (2)$$

Rearranging gives

$$NA_N = 1 + (N + 1)A_{N-1} \quad \text{for } N > 1 \quad (3)$$

Now divide the equation by  $N(N + 1)$  to get:

$$\frac{A_N}{N + 1} = \frac{1}{N(N + 1)} + \frac{A_{N-1}}{N} \quad \text{for } N > 1 \quad (4)$$

Iterating gives:

$$\frac{A_N}{N + 1} = \sum_{1 \leq j \leq N} \frac{1}{j(j + 1)} + \frac{A_1}{2} \quad \text{for } N > 1 \quad (5)$$

and with  $A_1 = 1$  from the original recurrence we have:

$$\frac{A_N}{N + 1} = \sum_{1 \leq j \leq N} \frac{1}{j(j + 1)} \quad \text{for } N > 1 \quad (6)$$

Noticing that  $\frac{1}{j(j+1)} = \frac{1}{j} - \frac{1}{j+1}$ , we write this sum as:

$$\sum_{1 \leq j \leq N} \frac{1}{j(j + 1)} = \sum_{1 \leq j \leq N} \left( \frac{1}{j} - \frac{1}{j + 1} \right) = 1 - \frac{1}{N + 1} = \frac{N}{N + 1} \quad (7)$$

and so:

$$A_N = N \quad \text{for } N > 1 \quad (8)$$

Since  $A_1 = 1$  from the original recurrence, we have:

$$A_N = N \quad \text{for } N > 0 \quad (9)$$

The way you've set up the recurrence, adding  $2/N$  is already captured by  $j = 1$  and  $j = N-1$ . Without this additive term you should be able to get  $E_N = (N+1)/3$ .

Equality holds whenever quicksort is called on an array of size 1. For  $N > 1$ , one of the two recursive calls will be on an array of size 1 if the pivot element is either the second largest or the second smallest element, which occurs with probability  $2/N$ . If  $N = 1$ , then the first and only call to quicksort is one on an array of size 1, so in that case equality holds exactly once. We can see a recursive relationship for the number of expected calls with equality. Let's denote  $E_N$  the number of expected equality calls for an array of size  $N$ . Then we have that  $E_1 = 1, E_0 = 0$  since for an array of size 1 there is precisely 1 call with equality (the initial call), and that:

$$E_N = \frac{2}{N} + \frac{1}{N} \sum_{1 \leq j \leq N} (E_{j-1} + E_{N-j}) \quad \text{for } N > 1 \quad (10)$$

We can solve this recurrence in the usual way – first multiply by  $N$  and note the symmetry of the sum term:

$$NE_N = 2 + 2 \sum_{1 \leq j \leq N} E_{j-1} \quad \text{for } N > 1 \quad (11)$$

Next subtract the same equation for  $N - 1$ :

$$NE_N - (N - 1)E_{N-1} = 2E_{N-1} \quad \text{for } N > 1 \quad (12)$$

Rearrange terms:

$$NE_N = (N + 1)E_{N-1} \quad \text{for } N > 1 \quad (13)$$

Divide by  $N(N + 1)$

$$\frac{E_N}{N + 1} = \frac{E_{N-1}}{N} \quad \text{for } N > 1 \quad (14)$$

Iterating gives the base case:

$$\frac{E_N}{N + 1} = \frac{E_1}{2} \quad \text{for } N > 1 \quad (15)$$

and substituting in  $E_1 = 1$  gives us:

$$E_N = \frac{N + 1}{2} \quad \text{for } N > 1 \quad (16)$$

Note that for  $N = 1$ , this formula also gives  $E_1 = 1$  so we can write generally:

$$E_N = \frac{N + 1}{2} \quad \text{for } N > 0 \quad (17)$$