

PS1-Q1

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5/5

How many subarrays of size 2 or less are encountered, on the average, when sorting a random file of size N with quicksort? When we have an n -element subarray called, the recurrence representing this situation is

$$A_N = \frac{2}{N} \sum_{1 \leq j \leq N} A_{j-1}$$

since it's the same as the recurrences we've solved in class/question 1.14, but without "adding anything" (as with the "1 +" in the second one) until we hit the base case (when size = n). Thus we have

$$NA_N = 2 \sum_{1 \leq j \leq N} A_{j-1}$$

$$NA_N - (N-1)A_{N-1} = 2A_{N-1}$$

$$A_N = \frac{N+1}{N} A_{N-1}.$$

Now, since we are looking for subarrays of size 2 or less, we will look for subarrays of size 2, 1, and 0 separately and add the expected numbers.

Size 2:

Telescoping, we see that

$$A_N = \frac{N+1}{4} A_3.$$

We note that for A_0 and A_1 we have 0 calls with subarrays size 2, for A_2 we have 1 (the whole thing), and for A_3 we have $2/3$ expected (choosing the first or last element gives a subarray of size 0 and size 2 called, one on either size, thus $1 \cdot 1/3 + 1 \cdot 1/3$, but the middle element gives two 1-element subarrays called = $0 \cdot 1/3$, thus $2/3$ total). (Note that since earlier in the problem we subtract the $N-1$ case, we must use the upper of two legitimate base cases. Normally $N=2$ would work, but we must use 3 for it to work here for the stated reason.) Thus

$$A_N = \frac{2(N+1)}{3 \cdot 4} = \frac{N+1}{6}.$$

Size 1:

Telescoping, we also see that

$$A_N = \frac{N+1}{3} A_2.$$

We note that for A_0 we have 0 calls with hi=lo, for A_1 we have 1, for A_2 we have 1 (one element is chosen as the pivot). (Note that since earlier in the problem we subtract the $N - 1$ case, we must use the upper of two legitimate base cases. Normally $N=1$ would work, but we must use 2 for it to work here for the stated reason.) Thus

$$A_N = \frac{N+1}{3}.$$

Size 0:

Telescoping, we also see that

$$A_N = \frac{N+1}{4} A_3.$$

We note that for A_0 we have 1 size-0 subarray called. For A_1 we have 0 as well, since Quick terminates after calling either a 1 or a 0 element array. For 2 we have 1, since either element it chooses will result in one 0-element subarray called on one side, and a 1-element subarray on the other. For 3 we have 4/3, because we have 2/3 for either end element (1 for the size 0 subarray, 1 for the size 2 subarray), and 0 for the middle element. Thus we have

$$A_N = \frac{4(N+1)}{3 * 4} = \frac{N+1}{3}.$$

Thus (except for some very small values of N , where our base cases were above them, but in general) the expected number of calls to arrays less than or equal to size 2 is $\frac{N+1}{3} + \frac{N+1}{6} + \frac{N+1}{3} = \frac{5(N+1)}{6}$.

In the future please list out the small cases explicitly i.e. 1, 1, 3